

An Overview on performance Analysis of AODV Protocol under Black Hole Attack

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ABSTRACT

Portable Ad-hoc networks are an assortment of versatile hosts that speak with one another with no foundation. Because of safety weaknesses of the steering conventions, remote impromptu organizations might be unprotected against assaults by the vindictive hubs. One of these assaults is the Black Hole Attack against network honesty retaining all information bundles in the organization. Since the information parcels don't arrive at the objective hub by virtue of this assault, information misfortune will happen. In this paper are doing reproduction investigation of organization under dark opening assault and do examination with the organization without assault chipping away at AODV convention utilizing different execution measurements like throughput, PDF and End to End delay in three distinct situations.

Keywords : Catchphrases Ad hoc network, dark opening , AODV, MANET, PDR, RREQ, RREP

I. INTRODUCTION

Mobile Ad-Hoc network is the system of mobile computer nodes or stations that are not physically wired. The main advantage of this is communicating with rest of the world while being mobile. The disadvantages are their limited bandwidth, memory, processing capabilities and open medium. Two basic system models are fixed backbone wireless system and Wireless Mobile Ad hoc Network (MANET). An ad hoc network is a collection of nodes that do not rely on a predefined infrastructure to keep the network connected. So the functioning of ad hoc networks is dependent on the trust and co-operation between nodes. Nodes help each other in conveying information about the topology of the network and share the responsibility of managing the network. Hence in addition to acting as hosts, each mobile

node does the function of routing and relaying messages for other mobile nodes [1]. In these networks, besides acting as a host, each node also acts as a router and forwards packets to the correct node in the network once a route is established. In this paper we will study the ad hoc network with and without the black hole attack using performance metrics PDF, Throughput and End to End delay. Irrespective organization is the organization of versatile PC hubs or stations that are not truly wired. The principle benefit of this is speaking with rest of the world while being versatile. The impediments are their restricted data transfer capacity, memory, preparing abilities and open medium. Two fundamental framework models are fixed spine remote framework and Wireless Mobile Ad hoc Network (MANET). A specially appointed organization is an assortment of hubs that don't

depend on a predefined framework to keep the organization associated. So the working of specially appointed organizations is subject to the trust and co-activity between hubs. Hubs help each other in passing on data about the geography of the organization and offer the obligation of dealing with the organization. Thus as well as going about as hosts, every versatile hub does the capacity of steering and transferring messages for other portable hubs [1]. In these organizations, other than going about as a host, every hub additionally goes about as a switch and advances bundles to the right hub in the organization once a course is set up. In this paper we will consider the impromptu organization with and without the dark opening assault utilizing execution measurements PDF, Throughput.

AODV

The AODV [2,3] routing protocol is a reactive routing protocol therefore, routes are determined only when needed.

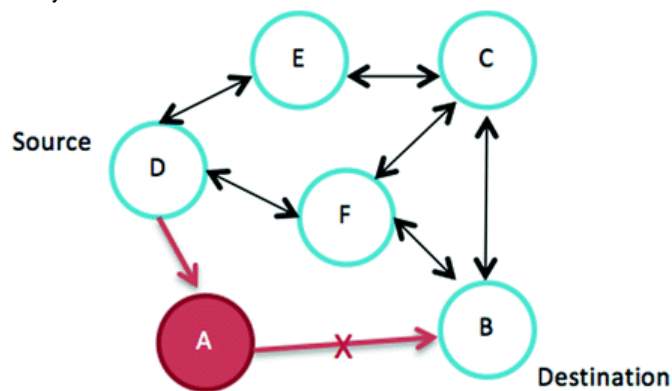


Figure 1.1

This figure 1 shows the message exchanges of the AODV protocol. Hello messages may be used to detect and monitor links to neighbors. If Hello messages are used, each active node periodically broadcasts a Hello message that all its neighbors receive. Because nodes

Periodically send Hello messages, if a node fails to receive several Hello messages from a neighbor, a link break is detected.

Source. As the RREP propagates, each intermediate node creates a route to the destination. When the source receives the RREP, it records the route to the destination and can begin sending data. If multiple RREPs are received by the source, the route with the shortest hop count is chosen. As data flow from the source to the destination, each node along the route updates the timers associated with the routes to the source and destination, maintaining the routes in the routing table. If a route is not used for some period of time, a node cannot be sure whether the route is still valid; consequently, the node removes the route from its routing table. If data is flowing and a link break is detected, a Route Error (RERR) is sent to the source of the data in a hop-by-hop fashion. As the RERR propagates towards the source, each intermediate node invalidates routes to any unreachable destinations. When the source of the data receives the RERR, it invalidates the route and reinitiates route discovery if necessary.

BLACK HOLE ATTACK

Routing protocols are exposed to a variety of attacks. Black hole attack is one such attack in which a malicious node makes use of the vulnerabilities of the route discovery packets of the routing protocol to advertise itself as having the shortest path to the node whose packets it wants to intercept [4]. This attack aims at modifying the routing protocol so that traffic flows through a specific node controlled by the attacker. During the route discovery process, the source node sends route discovery packets to the intermediate nodes to find fresh path to the intended destination. Malicious nodes respond immediately to the source node as these nodes do not refer the routing table. The source node assumes that the route discovery process is complete, ignores other route reply messages from other nodes and selects the path through the malicious node to route the data packets. The malicious node does this by assigning a high sequence number to the reply packet. The attacker

now drops the received messages instead of relaying them as the protocol requires. Malicious nodes take over all routes by attacking all route request messages. Therefore the quantity of routing information available to other nodes is reduced. The malicious nodes are called black hole nodes. For example, source A wants to send packets to destination D, in figure1, source A initiates the route discovery process. Let M be the malicious node which has no fresh route to destination D. M claims to have the route to destination and sends join reply JREP packet to S. The reply from the malicious node reaches the source node earlier than the reply from the legitimate node, as the malicious node does not have to check its routing table as the other legitimate nodes. The source chooses the path provided by the malicious node and the data packets are dropped. The malicious node forms a black hole in the network and this problem is called black hole problem.

II. SIMULATION ENVIRONMENT

We have executed [5] Black Hole Attack utilizing NS- 2.34 in AODV convention by adjusting the first convention and adding as another convention in NS2. X graph is utilized for plotting the outcome in type of chart in NS2.

The reproduction boundaries are displayed beneath.

Reproduction Area	800x800
No. of Nodes	20
Malicious Nodes	1,2...5
Correspondence Traffic	CBR
Reproduction Duration	10
Speed of the Node	20,50.....150 m/s
Parcel Size	512

We partition our work in three situations, in the first Situation we keep the speed of hubs and absolute number of hubs steady and we change the no of noxious hubs in the organization.

In the second situation we keep the quantity of hubs and speed steady, for correlation we add one noxious hub in the organization and think about the output as throughput, Pdf, start to finish delay with the first AODV convention organization.

In the third situation we keep the all-out nope. of hubs steady and one vindictive hub in the arrange and fluctuate the speed of the hubs and afterward think about the outcome again as throughput, pdf, start to finish delay with the first AODV convention organization.

The exhibition measurements expressed above are characterized as

Throughput: Throughput is the normal pace of effective message conveyance over a correspondence channel.

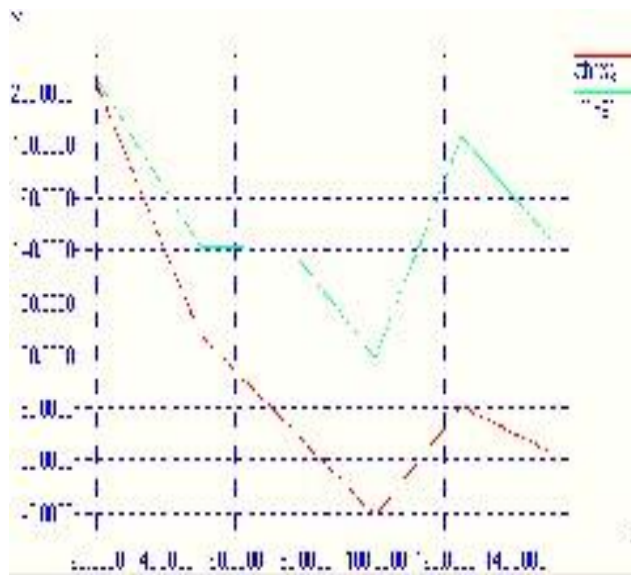


Figure 4.1 Throughput

2) Packet Delivery Ratio: The proportion between the quantity of parcels began by the "application layer" CBR sources and the quantity of bundles got by the CBR sink at the last objective.

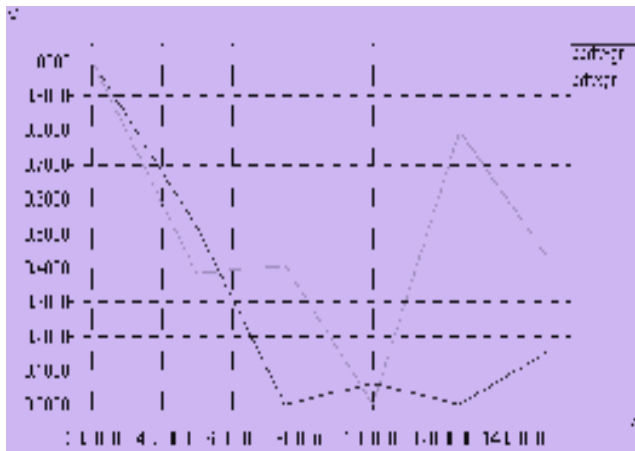


Figure 4.1 PDR

bundle conveyance proportion with the speed up hubs. With the correlation with the PDF of organization with no assault network with malignant assault have a lot of lower bundle conveyance proportion when the speed of the hub increments. For eg. The PDF in no assault network is 0.8 with hub versatility speed 125 m/s and the PDF in the organization with assault is 0.67 with a similar speed.

3) End to End Delay: Refers to the time taken for parcels to be communicated across an organization from source to objective.

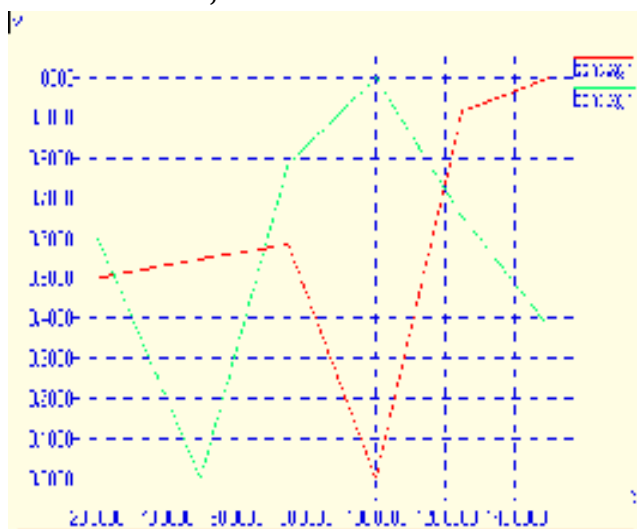


Figure 4.1 End to End Delay

E to E Delay with Black hole With no attack.

III. CONCLUSION

In this paper we have explored the introduction of off the cuff association under the dim opening attack and differentiated that and the association with no attack working using AODV coordinating show in three circumstances . The PDF and Throughput of the association has lessened certainly in all of the three circumstances and the End to End delay has extended again in all of the circumstances.

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